Analysis of Algorithms | Set 1 (Asymptotic Analysis)

***Why performance analysis?***  
There are many important things that should be taken care of, like user friendliness, modularity, security, maintainability, etc. Why to worry about performance?  
The answer to this is simple, we can have all the above things only if we have performance. So performance is like currency through which we can buy all the above things. Another reason for studying performance is – speed is fun!  
To summarize, performance == scale. Imagine a text editor that can load 1000 pages, but can spell check 1 page per minute OR an image editor that takes 1 hour to rotate your image 90 degrees left OR … you get it. If a software feature can not cope with the scale of tasks users need to perform – it is as good as dead.

***Given two algorithms for a task, how do we find out which one is better?***  
One naive way of doing this is – implement both the algorithms and run the two programs on your computer for different inputs and see which one takes less time. There are many problems with this approach for analysis of algorithms.  
1) It might be possible that for some inputs, first algorithm performs better than the second. And for some inputs second performs better.  
2) It might also be possible that for some inputs, first algorithm perform better on one machine and the second works better on other machine for some other inputs.

[Asymptotic Analysis](http://en.wikipedia.org/wiki/Asymptotic_analysis) is the big idea that handles above issues in analyzing algorithms. In Asymptotic Analysis, we evaluate the performance of an algorithm in terms of input size (we don’t measure the actual running time). We calculate, how the time (or space) taken by an algorithm increases with the input size.  
For example, let us consider the search problem (searching a given item) in a sorted array. One way to search is Linear Search (order of growth is linear) and the other way is Binary Search (order of growth is logarithmic). To understand how Asymptotic Analysis solves the above mentioned problems in analyzing algorithms, let us say we run the Linear Search on a fast computer **A** and Binary Search on a slow computer **B** and we pick the constant values for the two computers so that it tells us exactly how long it takes for the given machine to perform the search in seconds. Let’s say the constant for **A** is 0.2 and the constant for **B** is 1000 which means that A is 5000 times more powerful than B. For small values of input array size n, the fast computer may take less time. But, after a certain value of input array size, the Binary Search will definitely start taking less time compared to the Linear Search even though the Binary Search is being run on a slow machine. The reason is the order of growth of Binary Search with respect to input size is logarithmic while the order of growth of Linear Search is linear. So the machine dependent constants can always be ignored after a certain value of input size.  
Here are some running times for this example:  
**Linear Search running time in seconds on A**: 0.2 \* n  
**Binary Search running time in seconds on B**: 1000\*log(n)

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|n | Running time on A | Running time on B |

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|10 | 2 sec | ~ 1 h |

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|100 | 20 sec | ~ 1.8 h |

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|10^6 | ~ 55.5 h | ~ 5.5 h |

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|10^9 | ~ 6.3 years | ~ 8.3 h |

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***Does Asymptotic Analysis always work?***  
Asymptotic Analysis is not perfect, but that’s the best way available for analyzing algorithms. For example, say there are two sorting algorithms that take 1000nLogn and 2nLogn time respectively on a machine. Both of these algorithms are asymptotically same (order of growth is nLogn). So, With Asymptotic Analysis, we can’t judge which one is better as we ignore constants in Asymptotic Analysis.  
Also, in Asymptotic analysis, we always talk about input sizes larger than a constant value. It might be possible that those large inputs are never given to your software and an algorithm which is asymptotically slower, always performs better for your particular situation. So, you may end up choosing an algorithm that is Asymptotically slower but faster for your software.

we will take an example of Linear Search and analyze it using Asymptotic analysis.

We can have three cases to analyze an algorithm:  
1) Worst Case  
2) Average Case  
3) Best Case

Let us consider the following implementation of Linear Search.

* C++
* C
* Java
* Python3
* C#
* PHP

filter\_none

edit

play\_arrow

brightness\_4

|  |
| --- |
| // C++ implementation of the approach  #include <bits/stdc++.h>  using namespace std;    // Linearly search x in arr[].  // If x is present then return the index,  // otherwise return -1  int search(int arr[], int n, int x)  {      int i;      for (i=0; i<n; i++)      {      if (arr[i] == x)          return i;      }      return -1;  }    // Driver Code  int main()  {      int arr[] = {1, 10, 30, 15};      int x = 30;      int n = sizeof(arr)/sizeof(arr[0]);      cout << x << " is present at index "                << search(arr, n, x);        getchar();      return 0;  }    // This code is contributed  // by Akanksha Rai |

**Output:**

30 is present at index 2

**Worst Case Analysis (Usually Done)**  
In the worst case analysis, we calculate upper bound on running time of an algorithm. We must know the case that causes maximum number of operations to be executed. For Linear Search, the worst case happens when the element to be searched (x in the above code) is not present in the array. When x is not present, the search() functions compares it with all the elements of arr[] one by one. Therefore, the worst case time complexity of linear search would be Θ(n).

**Average Case Analysis (Sometimes done)**  
In average case analysis, we take all possible inputs and calculate computing time for all of the inputs. Sum all the calculated values and divide the sum by total number of inputs. We must know (or predict) distribution of cases. For the linear search problem, let us assume that all cases are [uniformly distributed](http://en.wikipedia.org/wiki/Uniform_distribution_%28discrete%29) (including the case of x not being present in array). So we sum all the cases and divide the sum by (n+1). Following is the value of average case time complexity.

Average Case Time = [analysis1](https://media.geeksforgeeks.org/wp-content/cdn-uploads/analysis1.png)

= [analysis2](https://media.geeksforgeeks.org/wp-content/cdn-uploads/analysis2.png)

= Θ(n)

**Best Case Analysis (Bogus)**  
In the best case analysis, we calculate lower bound on running time of an algorithm. We must know the case that causes minimum number of operations to be executed. In the linear search problem, the best case occurs when x is present at the first location. The number of operations in the best case is constant (not dependent on n). So time complexity in the best case would be Θ(1)  
Most of the times, we do worst case analysis to analyze algorithms. In the worst analysis, we guarantee an upper bound on the running time of an algorithm which is good information.  
The average case analysis is not easy to do in most of the practical cases and it is rarely done. In the average case analysis, we must know (or predict) the mathematical distribution of all possible inputs.  
The Best Case analysis is bogus. Guaranteeing a lower bound on an algorithm doesn’t provide any information as in the worst case, an algorithm may take years to run.

For some algorithms, all the cases are asymptotically same, i.e., there are no worst and best cases. For example,[Merge Sort](http://en.wikipedia.org/wiki/Merge_sort). Merge Sort does Θ(nLogn) operations in all cases. Most of the other sorting algorithms have worst and best cases. For example, in the typical implementation of Quick Sort (where pivot is chosen as a corner element), the worst occurs when the input array is already sorted and the best occur when the pivot elements always divide array in two halves. For insertion sort, the worst case occurs when the array is reverse sorted and the best case occurs when the array is sorted in the same order as output.